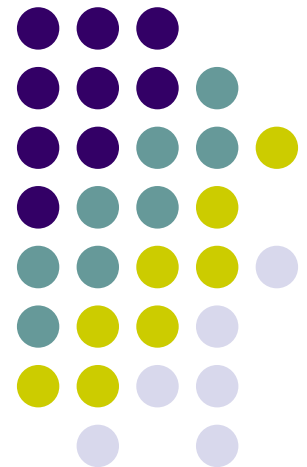


Joint Power Control and Link Scheduling in Wireless Networks for Throughput Optimization

Liqun Fu, Soung Chang Liew, Jianwei Huang

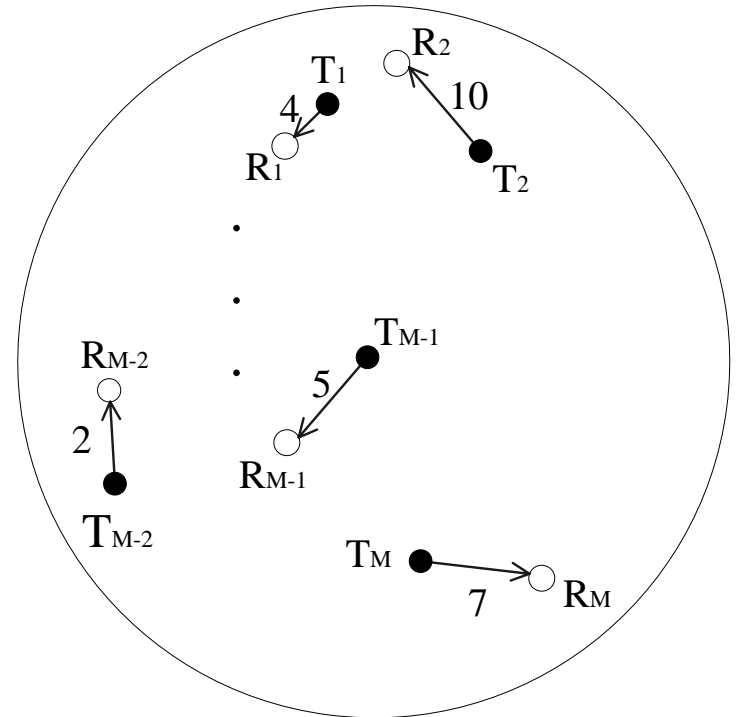
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Introduction



- Wireless scheduling
 - Single channel, TDMA network
 - Traffic demands
 - SINR constraints
- Objective
 - minimize the total number of time slots
- Power control
 - more links can be active simultaneously
- How to choose **the active links** in each time slot and the corresponding transmit **power** ?

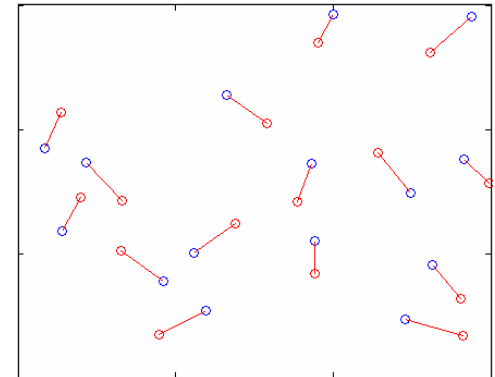




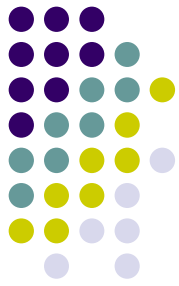
Challenge & Solution

- Challenge: too many feasible subset of links
 - Feasible subset of links
 - There exists non-negative power to satisfy SINR threshold

Number of links	15
Possible combinations	$2^{15} = 32,768$
Number of feasible subset of links	11949
Number of feasible subset of links (appear in optimal solution)	23



- Solution: Column Generation method



Column Generation Method

- Decompose the problem into two problems

Master Problem

- Consider only a small number of feasible subset of links



Pricing Problem

- Check whether the solution to the Master Problem is optimal
- Find a feasible subset of links which will further decrease the total number of time slots

- The Pricing Problem is difficult

The Pricing Problem



maximize: $\sum_{i=1}^M \omega_i y_i$

subject to: $\frac{p_i \cdot y_i \cdot G(T_i, R_i)}{N + \sum_{j=1}^M p_j \cdot y_j \cdot G(T_j, R_j)} \geq y_i \cdot \gamma_0$

variables: $p_i \geq 0$, if $y_i = 1$
 $p_i = 0$, if $y_i = 0$
 $y_i \in \{0, 1\}$



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**The subset of links
 \mathbf{y} is feasible**



$$\frac{1}{\rho(\mathbf{B}_y)} > \gamma_0$$

Power solution \mathbf{P}_y : the Perron eigenvector corresponding to $\rho(\mathbf{B}_y)$

The Pricing Problem



maximize: $\sum_{i=1}^M \omega_i y_i$

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Algorithms



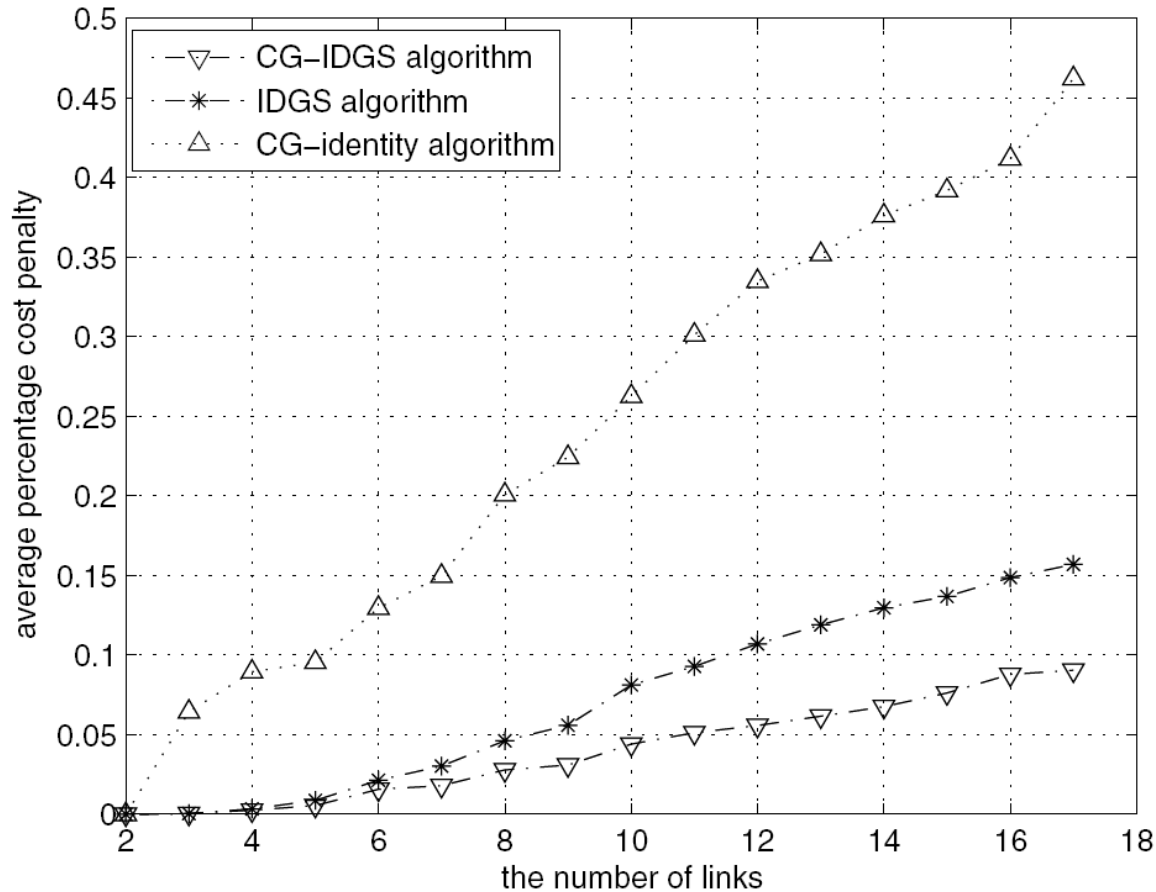
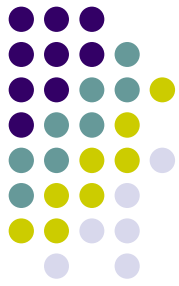
- Solve the Pricing Problem
 - Combined Sum Criterion Selection (CSCS)
- The Initial Feasible Subset
 - A simple Initial Solution
 - Each subset only contains one link: Identity
 - A good Initial Solution
 - Decrease the number of iterations
 - **Increasing Demand Greedy Scheduling (IDGS)**



Simulation Results

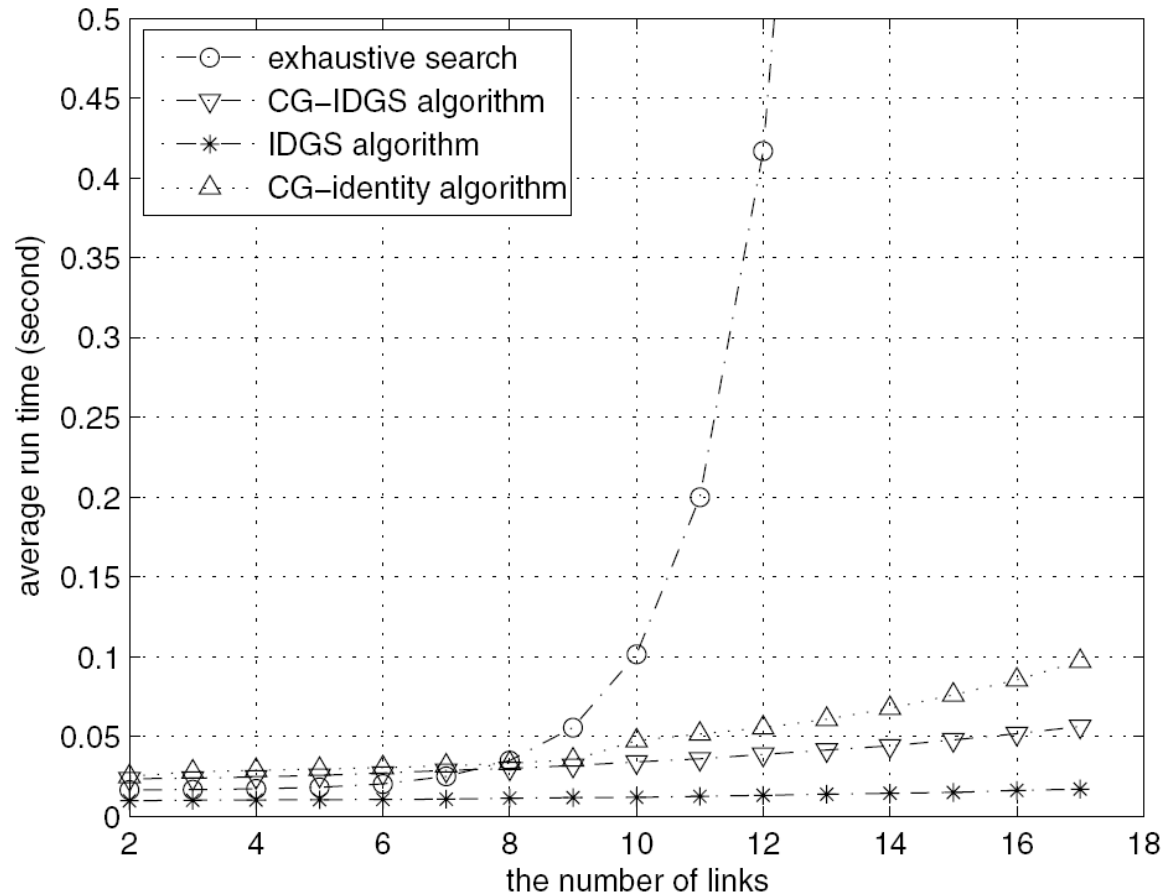
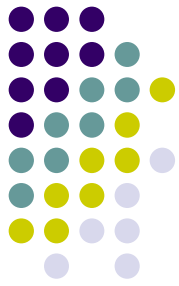
- Simulation setting:
 - Location: uniformly distributed in a square ($1km \times 1km$)
 - Path loss model: exponent 4.
 - Traffic requirement: discrete r.v. $[1, 3, 5, \dots, 19]$
 - Averaged over 1000 random networks.
- Performance comparison:
 - IDGS algorithm
 - CG-identity: Column Generation
 - Initial feasible subsets: each subset only consists one link
 - CG-IDGS: Column Generation
 - Initial feasible subsets: generated by IDGS algorithm

Simulation Result (1)



$$\frac{L - L_{opt}}{L_{opt}} \times 100\%$$

Simulation Result (2)



Conclusion & Extension Work



- Conclusion:
 - Propose Column Generation Method
 - Simplify the Pricing Problem without any loss.
 - Column generation + IDGS achieve great performance in short run time.
- Extension Work
 - A general wireless network: half duplex, non-simultaneously reception
 - Different SINR requirements at different receivers
 - The power is limited by a certain value.

Thanks!

